GT.M Update

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Empowering the Financial World



Overview

- Business
- Technical



Business









Monday signed a five-year framework agreement with the Electronic Health Solutions Company (EHS) to regulate the implementation of the e-health programme "Hakeem" at the national level.

The agreement — signed by Planning and International Cooperation Minister Imad Fakhoury and EHS Chairman Rami Farraj — aims at coordinating cooperation between EHS — the implementing party of Hakeem programme — and the concerned



Planning and International Cooperation Minister Imad Fakhoury and FHS Chairman

Technical



VistA on GT.M on Windows

Hats off to Christopher Edwards, Sam Habiel, and Cygwin



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Releases

• GT.M

- V6.3-000 major new enhancements
- V6.3-000A timely availability of fixes
- Plugins
 - Peer replication



GT.M Functionality

- Security Reduce "attack surface"
- Performance
- And more



Change database encryption password "on the fly"

- Database files vulnerable as a consequence of large data volumes and long lived keys
- Defense is to change keys periodically
 - Pre-V6.3-000: extract and load into new database
 - V6.3-000: Define two keys in configuration file, and run MUPIP REORG ENCRYPT to change encryption key while database is in use
- Can also encrypt unencrypted database files



Other Security Enhancements

- TLS key renegotiation
 - Protection against attacker who records a session for future
 - Effect with WRITE /TLS("renegotiate"[,,[tlsid][,,options]])
- Initialization vectors for database encryption different for each block
 - Better protection against attacks seeking to exploit large volumes of data
 - Automatic for databases created with V6.3-000
 - Zero IV retained for databases created with older releases changing the encryption key switches to using initialization vectors

Performance and Scalability

- Faster process termination
 - Especially when large numbers of processes and many regions both exist
- Option to allow more than 32Ki concurrent processes accessing a database region
- String expressions concatenation, \$[Z]ASCII(), \$[Z]EXTRACT(), \$[Z]PIECE(), and \$ZSUBSTR() on literals computed at compile time
 - In UTF-8 mode can result in BADCHAR warnings at compile time (errors at run time)
- Performance optimization for certain common UTF-8 mode use cases
- MUPIP JOURNAL commands have PARALLEL[=n] qualifier
- Enhanced replication throughput

And more...

- %PEEKBYNAME() gives symbolic access to internal statistics provided by \$ZPEEK()
 - LISTALL[%]PEEKBYNAME / LIST[%]PEEKBYNAME(.var) list all the fields, but sheer number is overwhelming
 - Fields not documented because internals can change from release to release, even if most may not work with GT.M support on specific needs
- MUPIP JOURNAL ROLLBACK FORWARD
 - Better / faster recovery from certain catastrophic operational failures, e.g., affecting multiple locations

Plugins

- Extend GT.M with useful functionality, but not part of core GT.M
 - Different release cycles from GT.M, and one another
- Usable "out of the box", but intended to be extensible
 - "Copy and extend" is a common programming model
- Released under the same license as GT.M
- Unmodified plugins supported as part of GT.M support, modified plugins not supported by FIS
- Existing plugins zlib, POSIX, GTMJI
- New Peer Replication



Serialization of Updates / Transactions

- Business logic that needs serialization
 - Order is required, e.g., each transaction on an account depends on the result of the preceding transaction on that account; a diagnosis must be made or at least suspected before a treatment can be prescribed
 - Unserialized updates cannot be aggregated
- Business logic that doesn't need serialization
 - Order can be relaxed, e.g., a balance inquiry on an account does not depend on the result of the previous balance inquiry (involves updates because each access to an account must be recorded)
 - Unserialized updates can be aggregated



BC / SI Replication vs. Peer Replication

- Centralized update serialization performs better than distributed update serialization
 - As long as central decision maker is not resource constrained (CPU, RAM, IO)
 - Systems like VistA (even EHS Jordan's Hakeem) are nowhere near limits of today's computer systems
 - BC / SI replication to distribute serialized updates from central decision maker
 - Unidirectional, implemented by GT.M, which maintains order, and restores it on recovery from a partitioning event
- Distributed updates can be closer to client, are more resilient to network outages, can use more economical computing resources, and can have greater elasticity with varying load
 - Peer replication to aggregate updates from multiple service points
 - Uni- or bidirectional, implemented by application code, which recovers aggregation (but not order) on recovery from a partitioning event

Sample configuration



Peer Replication



Peer Replication

- Implemented with pure application level code
- Triggers to capture updates to be replicated using peer replication
- Push / pull processes to monitor & replicate captured updates between instances



Peer Replication Schematic





Peer Replication Distribution

- Reference implementation of plugin is one routine
- Packaged with small sample application and scripts to demonstrate / test software
- Documented in readme and with help: mumps -run GTMPeerRepl —help



Questions and Answers



Contact – note new address and office landline!

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